Computability

Module 5

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1 Partial and Total Functions

2 Primitive Recursive functions

3 Recursive functions





- A Partial Function f from X to Y ($f: X \rightarrow Y$) is a rule which assigns to every element of X at most one element of Y.
- **Example:** if R denotes the set of all real numbers, the rule f from R to R given by $f(r) = +\sqrt{r}$; is a partial function since f(r) is not defined as a real number when r is negative.
- A Total Function f from X to Y is a rule which assigns to every element of X a unique element of Y.
- **Example:** The rule f from R to R given by f(r) = |r| is a total function since f(r) is defined for every real number r.
- We consider total functions f from X^k to X, where $X = \{0, 1, 2, 3, ...\}$ or $X = \{a, b\}^*$.
- We denote $\{0,1,2,\dots\}$ by N and $\{a,b\}$ by Σ .
- X^k is the set of all k-tuples of elements of X.



Partial and Total Functions



- For example, f(m, n) = m n defines a partial function from N to itself as f(m, n) is not defined when m n < 0.
- But g(m, n) = m + n defines a total function from N to itself.
- A partial or total function f from X^k to X is also called a function of k variables and denoted by $f(x_1, X_2, \ldots, X_k)$.
- For example, $f(x_1, x_2) = 2x_1 + x_2$ is a function of two variables: f(1, 2) = 4; 1 and 2 are called arguments and 4 is called a value.
- $g(w_1, w_2) = w_1w_2$ is a function of two variables $w_1, w_2 \in \Sigma^* : g(ab, aa) = abaa, ab, aa$ are called arguments and abaa is a value.





- The initial functions over N are given as:
 - **1** Zero function Z defined by Z(x) = 0
 - 2 Successor function S defined by S(x) = x + 1
 - 3 Projection function U_i^n defined by $U_i^n(x_1,...,x_n) = x_i$
 - 4 As $U_1^1(x) = x$ for every x in N. U_1^1 is simply the identity function. So U_i^n is also termed a generalized identity function.
- The initial functions over Σ are given as:
 - 1 nil(x) defined by $nil(x) = \land$
 - 2 cons a(x) defined by cons a(x) = ax
 - 3 cons b(x) defined by cons b(x) = bx





Example:

$$Z(7) = 0$$

 $S(4) = 5$
 $U_2^3\{2,5,7\} = 5$
 $nil(aabb) = \land$
 $cons\ a(aabb) = aaabb$
 $cons\ b(aabb) = baabb$



■ Composition of a function: If f_1, f_2, \ldots, f_k are partial functions of n variables and g is a partial function of k variables, then the composition of g with f_1, f_2, \ldots, f_k is a partial function of n variables defined by

$$g(f_1(x_1, x_2, ..., x_n), f_2(x_1, x_2, ..., x_n), ..., f_k(x_1, x_2, ..., x_n))$$

- The composition of g with f_1, f_2, \ldots, f_n is total when g, f_1, f_2, \ldots, f_n are total.
- Example: Let $f_1(x, y) = x + y$, $f_2(x, y) = 2x$, $f_3(x, y) = xy$ and g(x, y, z) = x + y + z be functions over N. Find the composition of g with f_1, f_2, f_3
- Solution: The composition of g with f_1, f_2, f_3 is given by $h(x, y) = g(f_1(x, y), f_2(x, y), f_3(x, y)) = (x + y) + (2x) + (xy) = x + y + 2x + xy$





- A function f(x) over N is defined by recursion if there exists a constant k (a natural number) and a function h(x, y) such that f(0) = k, f(n + 1) = h(n, f(n)
- **Example:** Define n! by recursion.
- Solution: Let f(0) = 1 and f(n+1) = h(n, f(n)), where h(x, y) = S(x) * y. So f(n) will be f(n) = h(n-1, f(n-1)) = S(n-1) * f(n-1) = n * f(n-1)
- A function f of n+1 variables is defined by recursion if there exists a function g of n variables, and a function h of n+2 variables, and f is defined as follows:

$$f(x_1, x_2, ..., x_n, 0) = g(x_1, x_2, ..., x_n)$$

 $f(x_1, x_2, ..., x_n, y + 1) =$
 $h(x_1, x_2, ..., x_n, y, f(x_1, x_2, ..., x_n, y))$





- \blacksquare A total function f over N is called primitive recursive
 - (i) if it is anyone of the three initial functions, or
 - (ii) if it can be obtained by applying composition and recursion a finite number of times to the set of initial functions.
- A total function is primitive recursive if it can be obtained by applying composition and recursion a finite number of times to primitive recursive functions f_1, f_2, \ldots, f_m . Each f_i is obtained by applying composition and recursion a finite number of times to initial functions.





■ A function f(x) over Σ is defined by recursion if there exists a 'constant' string $w \in \Sigma^*$ and functions $h_1(x, y)$ and $h_2(x, y)$ such that

$$f(\wedge) = w$$
$$f(ax) = h_1(x, f(x))$$
$$f(bx) = h_2(x, f(x))$$

 h_1 and h_2 may be functions in one variable.



■ A function $f(x_1, x_2, ..., x_n)$ over Σ is defined by recursion if there exists a function $g(x_1, x_2, ..., x_{n-1})$, $h_1(x_1, x_2, ..., x_{n+1})$, $h_2(x_1, x_2, ..., x_{n+1})$ such that

$$f(\wedge, x_2, \dots, x_n) = g(x_2, \dots, x_n)$$

$$f(ax_1, x_2, \dots, x_n) = h_1(x_1, x_2, \dots, x_n, f(x_1, x_2, \dots, x_n))$$

$$f(bx_1, x_2, \dots, x_n) = h_2(x_1, x_2, \dots, x_n, f(x_1, x_2, \dots, x_n))$$

 h_1 and h_2 may be functions of m variables, where m < n + 1.

A total function f over Σ is primitive recursive
 (i) if it is anyone of the three initial functions, or
 (ii) if it can be obtained by applying composition and recursion a finite number of times to the initial functions.





- Let $g(x_1, x_2, ..., x_n, y)$ be a total function over N. The function g is a regular function if there exists some natural number y_0 such that $g(x_1, x_2, ..., x_n, y_0) = 0$ for all values $x_1, x_2, ..., x_n \in N$.
- Example: g(x, y) = min(x, y) is a regular function since g(x, 0) = 0 for all $x \in N$.
- But f(x,y) = |x y| is not regular since f(x,y) = 0 only when x = y, and so we cannot find a fixed y such that f(x,y) = 0 for all x in N.

Recursive functions



- A function $f(x_1, x_2, ..., x_n)$ over N is defined from a total function $g(x_1, x_2, ..., x_n, y)$ by minimization if (a) $f(x_1, x_2, ..., x_n)$ is the least value of all y's such that $g(x_1, x_2, ..., x_n, y) = 0$ if it exists. The least value is denoted by $\mu_y(g(x_1, x_2, ..., x_n, y) = 0)$. (b) $f(x_1, x_2, ..., x_n)$ is undefined if there is no y such that $g(x_1, x_2, ..., x_n, y) = 0$
- In general, f is partial. But, if g is regular then f is total.
- A function is recursive if it can be obtained from the initial functions by a finite number of applications of composition, recursion and minimization over regular functions.
- A function is partial recursive if it can be obtained from the initial functions by a finite number of applications of composition, recursion and minimization.

Recursive functions



- **Example:** Show that f(x) = x/2 is a partial recursive function over N.
- Solution: Let g(x,y)=|2y-x| where 2y-x=0 for some y only when x is even. Let $f_1(x)=\mu_y(|2y-x|=0)$. Then $f_1(x)$ is defined only for even values of x and is equal to x/2. When x is odd, $f_1(x)$ is not defined $f_1(x)$ is partial recursive. As $f(x)=x/2=f_1(x)$ is a partial recursive function.
- **Exercise:** Show that $f(x, y) = x^2y^4 + 7xy^3 + 4y^5$ is primitive recursive.



Questions?

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Thank you.